#### A Micro-Mobility Management Scheme for Handover and Roaming debabala.swain@rediff\_nail.com **Debabala Swain** Dept.of Computer Science CUTM Bhubaneswar, India Siba Prasada Panigrahi siba panigrahy red fmail.com Dept of EEE GITA Bhubaneswar, India Prasanta Kumar Patra Dept. of Computer Science comuter@yahoo.co.in CET Bhubaneswar. India Abstract

Even though the PMIP provides mobility solutions, there are many issues of user identity, mobility context of users from a home network to the visiting network, the assignment of home address to a user terminal in a visiting network, identification of the user terminal's mobility, and identification of MPA and HA. In this paper, we propose a new uschangen with proxy mobile IPv4, as a mobility solution in networks. In this mechanism during mobile node access authentication, MPA exchanges registration messages with the HA (Home Agent) to set up a suitable routing and tunneling for packets from/to the MN. In this mechanism during network, this is then passed to the AAA (Authentication Authorization and Accranting) unver, and the authentication server checks the realm and does start authentication procedure at use time of initialing authorizing module of the mobile terminal. It also initiates the mobility extension module, where the AAA server initiates MPA of the access network, which also forms the AAA server of the home network with information on the mobility extensions and request of the mobility parameters of the user terminal. The home AAA server interacts with the HA and collects mobile node parameters, as well as sending back details as a noly request to the visiting AAA server. After the mobility context transfer, the MPA conducts a publity registration to the HA for that particular mobile node. Later in this paper, we will provide sequence of message exchanges during a mobility session of a user mobile node during half and request of the mosile during a mobility session of a user mobile node during half and request of the sequence of a user mobile node during half and request of the sequence of a user mobile node during half and request of the sequence of a user mobile node during half and request to the visiting AAA server.

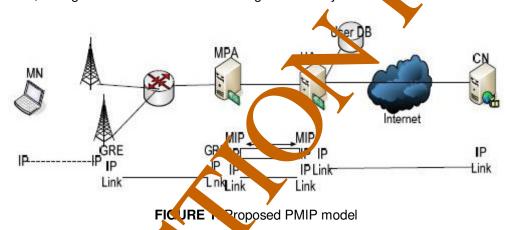
Keywords: Handover, Roaming, Mobility Management

## 1. INTPOD CHON

The mobility mana oppent in the access networks is provided by the mobile IP for the seamless continue of the sendees during handover and roaming. The demands for accessing services at high data haves while on the move, anyplace and anytime, resulted in numerous research efforts to integrate hereogeneous wireless and mobile networks. However, when the handover happens, the contention-based medium access mechanism which is mainly used in WLAN is involved and introduces unbounded transmission delay due to idle time periods and herensmission because of collision during the handover. If this technique is expanded to use in a micro relular network such as connected WLAN micro-cells, contention-based mechanism, therefore, should not be used to handle the MT's handover, especially for vehicular users who have access point every few seconds [1]. IP Mobility management protocols are divided into two kinds of category: host-based and network-based mobility protocol. Issues and challenges in

mobility management identified and discussed in [2]. Recently, a unified IP Multimedia Subsystem (IMS) authentication architecture that extends the scope of IMS by allowing it to offer users different IMS-based services even beyond their own domain has been proposed in [3] and all these research activities resulted in various heterogeneous architectures where the interworking was performed at different levels in the network. Also, integration at the UMAS radio access level for seamless session continuity proposed in [4]. But, proposed integration is a technology specific solution. However, in this article, we evaluate micro mobility.

The Proxy Mobile IP (PMIP) [5] solution based on Mobile IP approach; kandle mobility management inside access networks. Therefore network entities will require more napability than in the standard Mobile IP. The Foreign Agent is no longer capable to handle be mobility management in this new scenario, so we need to enhance its capabilities with the Mobily Proxy mechanism. This new entity called Mobile Proxy Agent replaces Foreign Agent in the visiting network. It also handles mobility registration with the Home Agent. This change is most significant since the Mobile Node now lies outside the mobility registration procedure. In hot, Mobile Node is not aware of its movement, access networks deceives the host to believe that his stationary in its Home Network. Since the Mobile Node does not need either poverient detection or agent registration, the agent advertisements are no longer necessary.



This paper addresses some of the requirements and features to be satisfied for PMIP to provide mobility management:

- Support Unmodified Hosts the noted above, the protocol supports mobility to nodes that does not have carebility of mobility.
- Air link consumption: would related signaling over the air-link is eliminated. Considering that Network Address Translation (NAT) is ubiquitous in IPv4 networks, a mobile node needs to send keep alive at short intervals to properly maintain NAT states. This can be performed by the MPA in the network which does not consume any air-link bandwidth. The Agent everticement is also eliminated in the protocol.
- Support the deterogeneous Wireless Link Network: One aspect is how to adopt the scheme to an access pennology. Since Proxy Mobile IPv4 is based on a heterogeneous mobility processl, it can be used for any type of access network.
- The one aspect is how to support mobility across different access technologies. As long as the MPA can use the same NAI to identify the MN for various access networks, roaming netween them is possible.

Support the IPv4 and IPv6: As IPv6 increases in popularity, the host will likely be dual stack.

# 2. PLOPOSED SOLUTION FOR PMIP WITH INTEGRATED AAA ARCHITECTURE OF THE 3GPP AND WIRELESS NETWORKS

In this new mechanism, mobility registration of a user terminal is performed by visiting access petworks and a home access network. The user terminal does general authentication by visiting

access networks with the help of an EAP (Extensive Authentication Protocol) mechanism. The visiting access networks receive the authentication request from a user terminal through the NAS or AP of the network. The AAA server of visiting network and home networks are modified so they can communicate with the HA and MPA of their respective networks. New mobility extensions are developed in AAA server to support mobility management, which adds to its present services. These extensions provide mobility context transfer from home access networks. registering the user terminal for mobility at the time of authentication. The visiting network prates authentication and the mobility extension method whenever it receives a request free the two or AP of the access network. During initiation of mobility extensions, the AAA mobility extension process collects data when NAS/AP requests authentication. The AAA mobility rogess sends mobility user details request to the home network and the AAA server of the terminal with newly specified attributes of proxy mobile IP. The Home AAA server does receive a requestor the mobility user details request as well as the authentication. The home AAA set distinguishes a proxy mobile IP packet from other codes and attributes of the received packet. If the packets need to be a proxy from an intermediate AAA server, then that server adds the proxy attribute to the received packet and sends it to the destination AAA server. If ever the use terminal belongs to the current network, then the AAA server sends a mobility regised tion request to HA.

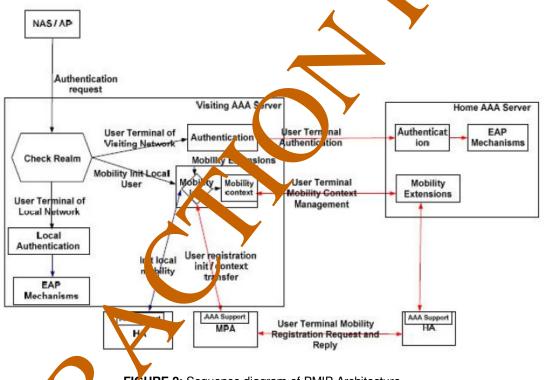


FIGURE 2: Sequence diagram of PMIP Architecture

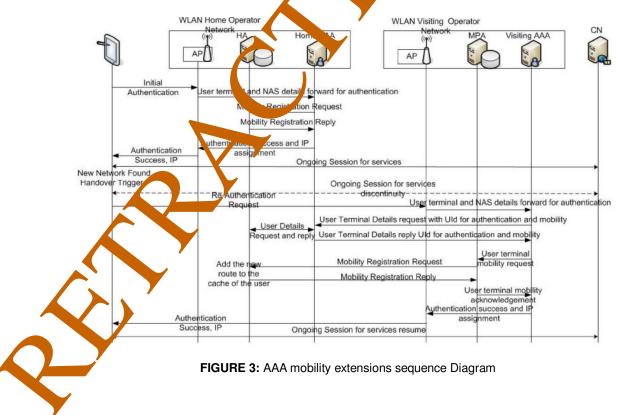
After receiving the equest for mobility user details packet from the visiting AAA server, the home AAA server investigates any information available in the packet and collects user identity from the request packet of the HA of the access network. This packet contains details of user id and parameters. The HA receives a request, and with a user ID of request it extract the information of its SID, keys, home address and home agent address from the database of the HA. The HA then such back a reply to the AAA of home network with the above mentioned data. The AAA server receives a reply and processes the information, and sends back a reply message to the visiting AAA server receives a reply from home server and processes it, storing the rate of the user in a temporary database. After processing the reply message, the AAA server sends a mobility registration request to the MPA associated with that particular NAS or AP. This

request contains the details about user ID, SPI, keys, home address and home agent address. When the MPA receives the packet it starts the mobility registration of a user with details from the AAA server.

MPA initiates a mobility registration request of a user terminal with HA using details provided by visiting AAA server. Registration involves the user SPI and the shared key mechanises with the key available from the AAA server to the MPA. After successful registration of the user vith the HA, the MPA will modify the DHCP server configuration with the user terminal's ils. ese modifications contain details of MAC address and home address of the user in the DHeP server. After successful authentication of the user terminal it initiates a DHCP request to be address. The AP/NAS of the visiting network forwards the request to the DHCP server. When the MAC address of the user terminal modified, the DHCP server sends a reply to user's termine with its home address. The user terminal receives the reply and configures the IP access to the home address. Necessary modification has to be done by the visiting network to a conmodate the terminal with the ARP, etc. When the user terminal is in it home demain, the HA registers the terminal and sends the modified DHCP request to the DHCP server and technomledges the home AAA server of successful registration of a user terminal. The Prov, Mobile IP with the AAA server mobility architecture is shown in Figure 1.

#### 2.1 AAA Mobility Extensions for PMIP Integrated Architecture

In this section we describe the detailed architecture of AAA with mobility extensions to provide mobility management during user mobility in different account works and technology. In this process, the existing AAA architecture is modified to accommodate proxy mobile IP. In general, authentication information of users is passed through the acthenticator, and then this information is passed through the NAS or AP of the access new rks. In AAA server authenticates the access networks for the AP or NAS initially, and then processes the user authentication request depending on the realm of the user. In this new method, the mobility management of a user can be initiated during the authentication process. In this process, due to parallel operation of authentication and mobility management, the operation of a user during the handover and initial access can be reduced.

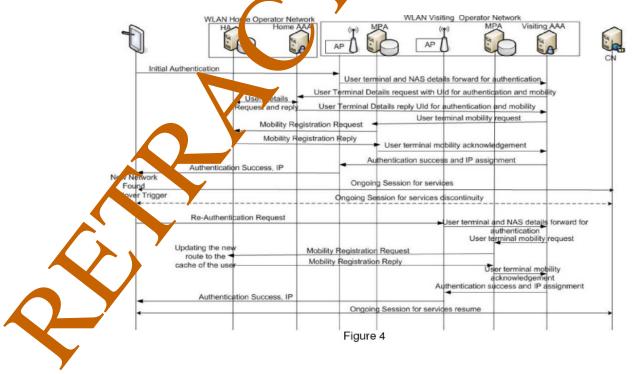


When there is an authentication request for a user terminal from an NAS or AP, the AAA server initiate authentication module and mobility modules, and processes the user's details by identifying the NAI of the user terminal request. From the NAS or AP request information, success MAC address of user terminal, NAS details are processed for further procedures. The AAA is modified, with new attributes and codes being added for supporting the PMIP modules. As mentioned previously in the proposed solution section with new extensions, the AAA of the home network can communicate with a visiting network, and can provide mobility context manufement. With these mobility extensions, the AAA server can communicate with the MPA of the access networks.

On the other side, the visiting AAA server communicates with the home network AAA arver, after receiving an authentication request using the mobility extensions, with user information being available from the authentication request from the user's terminal. The volume power sends a mobility user details request using ID and NAI of the authentication request to the nome AAA server. When the home network receives a request packet, the AAA server processes the information of the user from request. It then sends a request to one not with the new mobility extension, requesting details of the user. After receiving the request packet with home address, key, SPI and home agent address to the home AAA server. The home AAA server sends back a reply to the visiting AAA server with user details as the reply. After receiving this reply from the home AAA server, the visiting AAA server processes the information of ne user and sends a request for mobility registration request with new attributes to the APA Theorem PA receives the user terminal data, sent by the visiting AAA server, and temporarily stores it in a local database. The MPA, with available user information, starts registering with the HA. After registration request and reply message exchange with HA, the MPA sends reply of success of the failure of mobility registration of user to visiting AAA server. Figure 2 describes the AAA pobility architecture.

### 2.2 PMIP Operation With New Mobility Extensions in MPA and HA

MPA exchanges registration messages with the HA to set up a proper routing and tunneling packets from/to MN. The MN broadcasts messages containing an MN's Network Access Identifier (NAI) to request authentication/authentiation, and the AP transfers the request to the local AAA server (visiting AAA). If the MN is away from home, it is clear that the MN is out of the local authentication database.



Mobility extensions using MPA and HA

However, the local AAA server can use the NAI to identify the MN's Home Network, and then the authentication/authorization, along with mobility user details, will request a message to be transferred by the visiting AAA to the home AAA Server (AAAH) in the Home Network.

Along with the authenticating validation, the AAAH searches for information of the stored in the HA, containing MN's HA, NAI, and SPI. If the MN is back to its Home Network, the the local AAA server sends a message to the HA to deregister the MN instead of searches for the data. The MN's information will be transferred to the visiting AAA, which will deliver it to be MPA with the AP's MAC address included. Triggered by the AAA server, the MPA exchanges havesages with the HA to demand Mobility Registration and Tunneling.

After successful registration, the MPA sends a message to inform the DHCr server about the MN's arrival. It forces the DHCP server to update the configuration her with the Mobile Node information. Finally, the MPA informs the AAA visiting about the successful registration. The Authentication Accept message is sent to the NAS, granting networ access to the MN. After authentication success, the MN sends a Binding DHCPDISCOVER to cruest the IP address. This message is formatted as described by the DHCP protocol (the CIADDR field is filled with the MN's IP). By searching for information of the Mobile Node, in the configuration, the DHCP server replies with a DHCPOFFER message in which the YMDDR field is filled with the MN's Home Address and the default gateway address, being the MPA's. Next, the MN and DHCP server exchange the DHCPREQUEST and DHCPREPLY to complete this procedure. The MN is then ready to connect to the network with its Home Address.

## **3 MICRO MOBILITY**

In this scenario mobility is performed, same doministrative domain and same access technology, we have observed two sub menanes where the proposed architecture addresses this issue.

### 3.1 User Terminal Mobility in Home Administrative Domain on Same Access Technology

In this scenario access network has multiple APs, and user terminal moves from one AP to another AP. During initial authentication of user, AAA server does authenticate user and assists HA for mobility registration of user terminal. When user terminal senses other APs of access network and triggers the hand over with re-authentication procedure, upon receiving request from new AP, the AAA server sends a mobility registration request to MPA associated with AP. MPA and HA does the mobility registration of the user terminal and sends acknowledgement to AAA server. Upon successfully accessfully accessing terminal in HA, it provides access and home IP address of user terminal to AP for providing access to user terminal. The message exchange is shown in Figure 3.

**3.2 User Terminal Mobility in Visiting Administrative Domain on Same Access Technology** In this scenario a user terminal moves on same interface from one AP to another in visiting operator network. The user terminal is authenticated and registered in HA with the help of home AAA and visiting AAA servers. When user terminal identifies new AP it triggers the handover and doet re-autentication procedures. Upon receiving request from new AP visiting AAA identifies user from previous registration and sends mobility registration request to new MPA with previous details. MPA does register the user terminal with HA and sends acknowledgement to visiting AAA server to complete handover procedure, the whole procedure is shown in message sequence in here 4.

#### 3.3 Enhancing the Proposed Solution Using Network Selection Procedure for Seamless Mobility.

To enhance proposed architecture we used network selection procedures combined with this architecture to use context management between the networks. Using this process access networks can create mobility context even before user terminal does initiate access to visiting network. In this process user terminal can communicates with home network using precent connected network and negotiate best suitable network to connect during handow. After selecting best suitable network with the assistance of terminal, home network in the curtext transfer and creating mobility context with the future visiting network.

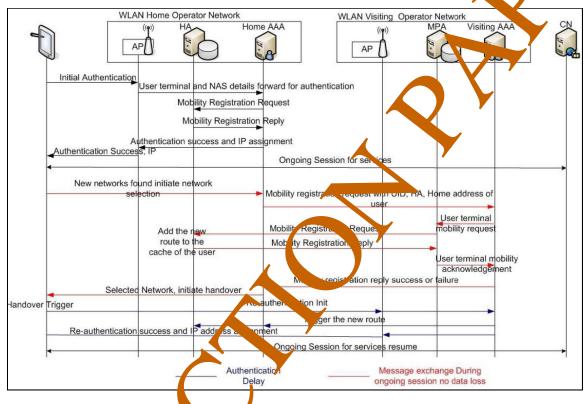


FIGURE Mobility nanagement using Micro mobility model

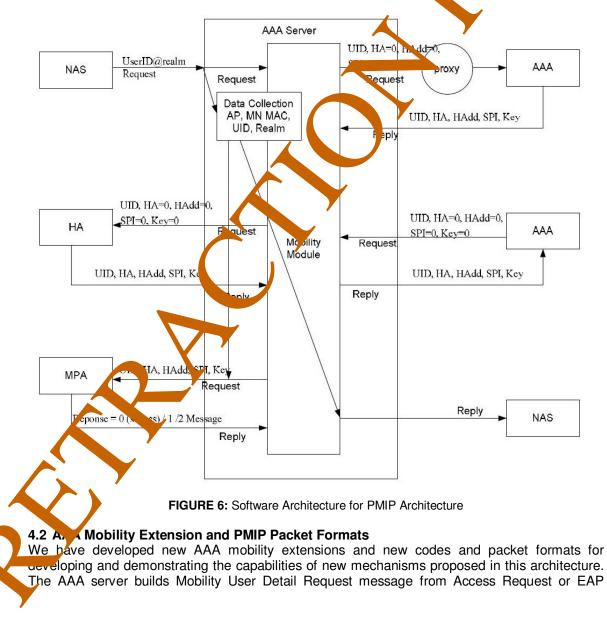
Using AAA mobility excertions proposed in this architecture AAA of home network sends a mobility registration request to waiting AAA server with UID of the terminal and mobility context details in the request. After receiving request from home AAA server, visiting AAA server collects data and sender registration request to MPA of visiting network. After receiving request for mobility registration may be accessed and initiates registration request to HA of home access network. After receiving registration user details of new route are cached in HA and MPA and a tunnel is established between them. When user or home AAA server does the handover triggering, the does update route upon receiving the RU (route update) request from home AAA server. In this way maintaining multiple tunnels with future visiting networks of user and biggering with the help of home AAA server seamless mobility is achieved.

The vinole message exchange sequence diagram is shown in Figure 5. During implementation of is procedure in a test bed we observed zero latency for multi homing handover and for how ontal handover we obtained small latency delay due to re-authentication procedure.

## 4 NEW PMIP AND AAA MOBILITY EXTENSION DEVELOPMENT AND TEST-BED SETUP

#### 4.1 PMIP and AAA Software Architecture

To implement proposed architecture we developed AAA server and PMIP in house using existing open source software. We have developed software architecture to implement mobility extensions for AAA server. In this architecture AAA server can receive a request from tAS or from another AAA server. From NAS it can receive authentication request and from AAA server is can receive mobility user detail request. Upon receiving mobility registration request, extensions model respond with the reply of user details. Software architecture of AAA as show in Figure 6, the AAA server can send requests and reply accordingly to incoming requests with an different components. For implementing the PMIP we used dynamics mobile IP architecture and modified to our requirements. We converted FA to an MPA, modified HA and MPA to accountly registration requests to HA upon request from AAA server and sending reply accordingly. In this architecture MPA can be account of from AAA server and sending reply accordingly. In this architecture MPA can be account of from requests or failure. New packet formats and codes are added in MPA and TA to implement the proposed architecture.



Request from the NAS or AP. Remark that intermediate AAA servers just pass through this step adding Proxy Attribute and forwarding the Request.

### 4.2.1 AAA Mobility User Detail Request Format:

Note: the codes and the attributes in this document are taken as reference these can be changed according to the IANA consideration; in this case we used available values for developing the prototype, we can change these values if there are any issues.

0 1 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 9 0 1 Code (1 byte) | Identifier(1B) | hvtes) Length (1 Authenticator (16 bytes Attributes ... Code: (1 byte) Mobility User Detail Request = 60. Identifier: (1 byte) number to match the Request Reply. Length: (2 bytes) length of the message, including Code, Identifier, Length, Authenticator, Attributes. In the case that there is only movility a bute, length = 350. Authenticator: The Authenticator field 50 by The most significant octet is transmitted first. This value is used to authenticate the ply from the RADIUS server, and is used in the password hiding algorithm. Attributes: Mobility Attribute: 0 1 3 3456789012345678901 0 1 2 3 4 5 6 7 8 0 1 2 Length (2 bytes) User's ID (1B) Type (1 byte User's ID (256 bytes) HA address (4 bytes) Home address (4 bytes) SPI (Lyte) Key (64 bytes) be = (1 byte) Mobility Request Attribute = 193. Let th (2 bytes) = Length of the message = 332. User ID: (256 bytes) extracted from the name of the user (ex: userID@realm). ddress: (4 bytes) Home Agent's IP address, filled with Zeros. Home Address: (4 bytes) Mobile Node's Home Address, filled with Zeros.

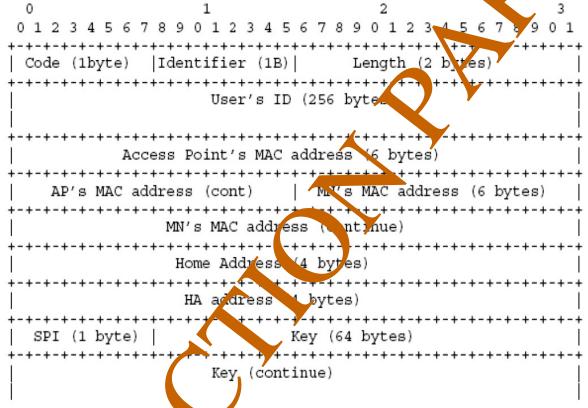
SPI: 1 byte, filled with Zeros.

Key: (64 bytes) public key of the HA, filled with Zeros. The AAAH will reply with a Mobility Response.

#### **HA/MPA Consultation**

If AAA home server receives Mobility user detail request from a visiting server, the AAAA sends message to HA to fill the information required in the Mobility Request Attribute (field's that are filled with Zeros). Remark that the AAAH sends the HA Consultation message only user detail Request; the Access Request forces the AAAH derevister the MN.

H1: Create a message from AAAH to the HA demanding for the necessary information



Code (1 byte) = HA\_Consultation\_request = 63. Identifier: (1 byte) num.c. match. Request/Response.

Length (2 bytes) = total length on the message = 343

AP and MN's MAC address: These fields are practically used in the message from AAA to MPA. In the message from AAA to HA, these fields are filled with Zeros, and the HA just ignores it. But these fields SHOU D appear in the HA Consultation Message to identify the format of messages AAA-HA and AA, mean is very useful since the HA and MPA in the same network are usually installed in the one server. This identification simplifies the treatment of message in the HA/MPP server.

#### Other fields a copied from the Mobility Attribute of the authentication request.

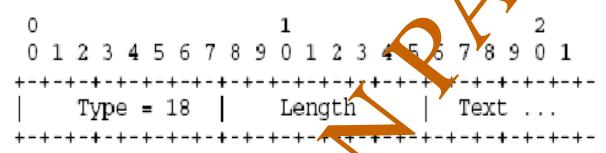
H2: HA looks for the required information in its database, save the AP and MAC address fields. If the information can't be found (this may be due to the modification of the administrator), HA will use this mase, so that the message will be left Zeros. That allows the AAAH to detect the failed.

H3: HA sends back the reply to the AAA after filling the request's required fields and setting Code = nA\_Consultation\_Response = 64.

H4: The AAAH replies the visiting AAA with a Mobility user detail Response, which is either an Accept or Reject message. The format of these messages is as same as the request, with different code and attributes.

If the message from HA is not filled with Zeros (successful verification), the AAAH reply to the AAAF with Mobility Accept message which is copied from the Mobility Request phose the Attributes filled by the data retrieved from HA. The Code field for this message is for the Mobility\_Accept = 61.

If the data from HA is filled with Zeros, the AAAH MUST reply the visiting AAA with a Mobility user detail Reject message, with Code = Mobility user Reject = 62. The Mobility Reject message doesn't contain the Mobility\_Attribute, and may include Reply-Massage Attribute which contains the error message shown to the user [6]:



#### Type: 18 for Reply-Message.

Length: length of the attribute, including Type and Length leld.

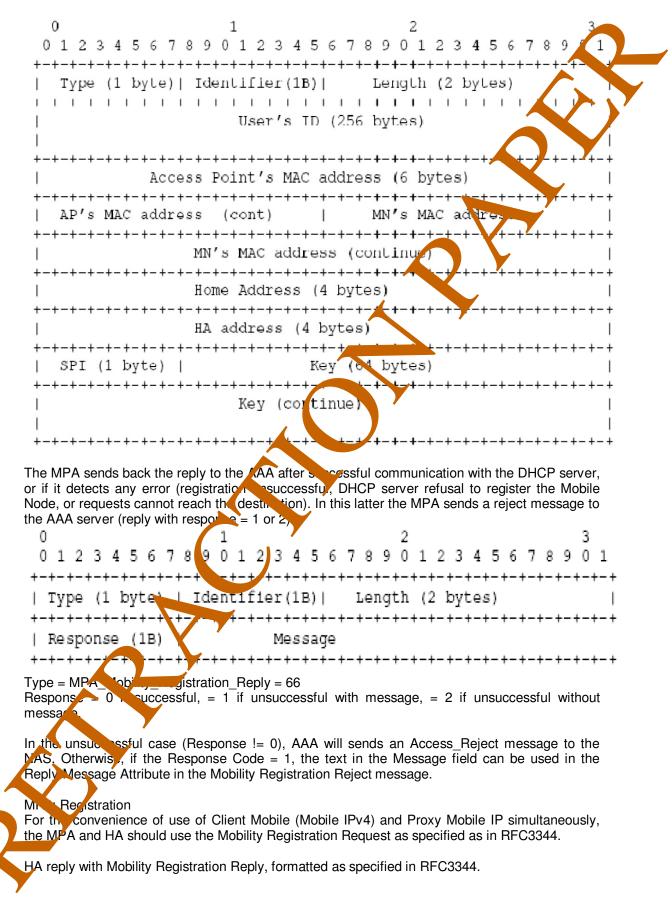
Text: The Text field is one or more octets, and its contents are implementation dependent. It is intended to be human readable, and MUST for affect operation of the protocol. If the registration failed, this field is filled with the message extracted from the MPA Mobility Registration Reply.

#### Mobility Registration

After receiving the Mobility Accept message, the visiting AAA makes MPA handle the Mobility Registration procedure. The M-A exchanges messages with HA and DHCP server, then informs visiting AAA about the result (success or failure). The Mobility registration Reject causes the AAAF to send the Reject message to the NAS and terminate the whole procedure.

MR1: Visiting AAA sends a MPA moonity Registration Request message to MPA: the format is as same as HA Consultation pessage;

Type (1 byte) = MPA\_Movinty\_Registration\_Request = 65. Identifier: (1 byte: number to match Request/Response. Length (2 bytes) = btal length of the message = 343. Other fields size / manufiles MAC address are copied from the Mobility Attributes of the Registration-Response. MR6: MPA Mobility registration Reply to visiting AAA:



MR4: MPA sends DHCP Mobility Registration Request to DHCP server: 0 2 1 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 | Type (1 byte) | Identifier (1 byte) 1E Length MAC address (6 bytes) 4 hvtes) MAC address (continue) | Home Address Home Address (continue) | Action (1B) | MA's MAC add MPA's MAC address (continue, (6 bytes) | MPA's MAC add | Type = DHCP Mobility Registration = 67 Identifier: match Request/Response Length = 21: length of the message, including the Type and Identifier fields MAC address: MN's MAC address Home Address = MN's Home Address. Action: (1 byte) = 0 - binding update: the DHC server updates its configuration file with the MN's new entry: MN's MAC address ---- MN's IP address --- Default Gateway = MPA's MAC address If Action = 1 - remove entry: cause the DHCP server to remove the MN's entry in its configuration file. This action is used in the Registration Revocation Procedure. As receiving the message from the MPA, the DHCP server to dates its configuration with the information supplied by the MPA. Since then, as soon as the DHCP server receives the (Binding) DHCPDSICOVER message from the MN, it will exchange the messages with the MN granting the MN keep its Home Address; also indicates the MPA as MNs control gateway. MR5: DHCP Mobility Redistration Reply to MPA The message is matted same as the reply from MPA to AAA 3 0 1 2 6,7890123456789012345678901 bvte) | Identifier(1B)| Length (2 bytes) Respinse (1B) Message = DHCP Mobility Registration Reply = 68 Length = length of the message including the Type and Identifier fields Response Code = 0: accept, = 1 reject with message, = 2 reject without message.

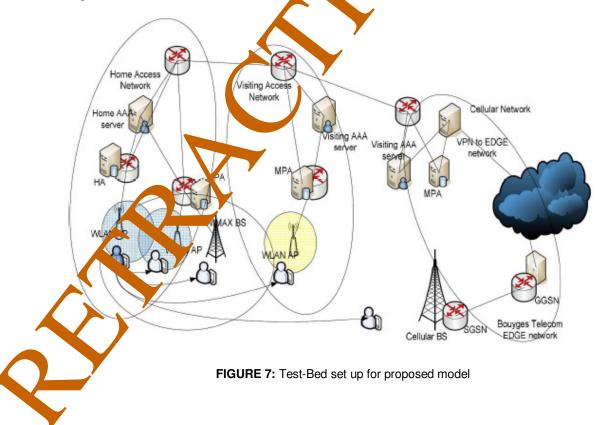
If the response Code is other than 0, the MPA MUST response with the AAA with MPA Mobility Registration Reply whose Response and Message fields copied from DHCP Mobility Registration Reply message.

Message: this message will be used in the response from MPA to AAA.

#### 4.3 Test bed Setup

This section describes testbed setup for implementing solutions proposed in this arc tectur. As mentioned earlier we have developed mobility extensions for AAA server using Free adius [1], and PMIP using parts of Dynamics mobile IP [8] with our implementation. The proposed testbed composed of 3GPP, WLAN and WIMAX Networks. We used Infinet's preWIMAX requipment, operating at the frequency of 5.4 GHZ for WIMAX network, WLAN access consists or Linksys WRT54 and Cisco AirNet AP350. The 3GPP network used in this care of FDGE network operated by the French network operator Bouyges Telecom courtesy of M/NO Transatel. The user terminal used in the testbed is DELL Latitude410 using Centring for diverent scenario we have used WIMAX and WLAN due to complexity of EDGE network acree have limited control of the production network provided by MVNO Transatel for us. To validate the solution we used basic testing like vertical handover where the client is equipped with multiple interfaces to validate the solution for multi homing scenarios. As we mentioned earner horizontal handover is limited for EDGE in our case because of the control of the access network.

The implemented scenario is shown in Figure 7. We have deployed a VPN with cellular network where authentication and data is routed through the cellular network to local testbed. The user terminal is configured with AT commands using PAP routile action. When user terminal dials for connection, authentication information of u er terminal is routed through GGSN to the local authentication in the test bed, the modified AAL server does the authentication and assigns the address using IP pool mechanism, and imparation mobility context is created for MPA and HA using AAA server.



We have deployed EAP authentication mechanism for authentication in WLAN and WIMAX networks. A user terminal tries to connect access networks using WPA supplicant [9], it is configured with user id with NAI and security mechanisms essential for EAP TLS mechanisms. Once the authentication is initiated in access networks, APs and BS sends authentication request to AAA servers, and AAA initiates authentication and mobility context for user terminal. A GUI is developed on terminal to maintain interfaces and control access management along different access networks.

# 5 **RESULTS**

As mentioned in architecture the AAA server does authentication and mobility in scallel when there is a request from the user terminal in test-bed. Using this test-bed we have achieved mobility of user with low latency and seamless mobility in some scenarios. Multi noming, horizontal handover and roaming is performed efficiently using this mechanism. Various scenarios of mobility have been tested using this test-bed. Deployment and extending to the new access networks and operator is very efficient as the modifications are made at network side without any client conscious. The modifications on the network side can be made with additional patches with existing deployments. For testing purposes we have ned experimental codes and attribute value pairs these can be extended to the vendor specific or using IANA status can be standardized.

We have observed an overall latency of the user terminal involved during roaming from one network to another on the same technology is around an overall N, 3.6 seconds for WIMAX, and 16 seconds for 3G networks, due to re-authentication and mobility management. In 3G networks we have observed high latency due to the lay on routing messages from bouyges telecom network to our test bed. For multi homing cenario we observed latency of 18 milliseconds as we have implemented multiple interface scenarios, where a user terminal connects to multiple networks and the management of the mobility is performed by triggering route update message in HA of the user to be next paragraph we have attached a log of our radius servers in home and visiting networks where the whole procedure is depicted. For more details about logs, ethereal results refer to paid

### 5.1 Comparison With Existing M bility bility

In this section different mobility protocols are compared with the available results and support for access networks with experimental results and simulation results. As we mentioned in last section we have built test bed to perform mobility in different interworking scenarios of mobility using mobility protocols. We take the following result of CMIP, HMIP and PMIP from the test bed and the result of HAWAII and Cello. The room other sources [11]. The Table 1 shows the different performance results for mobility protocols.

In the test for CMIP and AMIP, we take the result from mobility registration procedure only. Macro-mobility bandoff time is counted from moment that MN starts to the end of the handoff; and micro-mobility han off time is counted from the moment that we switch network connection (changing a cases usint) fact, using manual switch is little different from real time test in which the MN starts updoff if it moves out of the first access point's cover, since in the later case's handoff time depende largely on network scanning and selecting software used in the MN. We can observe that the CMIP has very high macro-mobility latency, which dues to Agent Discovery phase. We usnit that the test bed is so simple as compared to the architecture implemented H WAII and contular IP that the different networks are adjacent, and therefore cannot have a proper comparison among these protocols. The test is purely on latency issue, we don't count on packet loss and robustness. However, the result is persuading enough to prove the advantage of the prove t

Proxy Mobile IP is advantageous over other mobility protocols over security, since the information changed among the network entities with authenticated mechanism. More precisely, the advantage of PMIP over other protocols comes from the fact that the information exchanged in

registration procedure can be generated for each session, i.e., HA can generate necessary information used for each registration session. Hence, outside AAA authentication, no key is actually stored for mobility registration.

The proposed mechanism for mobility management in this paper is compatible and interoperable with the existing converging networks. We have studied different interworking methods to implement our solution for completing the seamless converging puzzle at the mobility management layer. We interrogated different interworking mechanisms such in Seauless Converged Communications Across Networks (SCCAN), Unlicensed Mobile Access (UNn), Interworking- Wireless LAN (I-WLAN), Media Independent Handover (MIH) Have 8/2.21. The proposed solution can be adapted in these mechanisms to provide seamless services at the mobility layer.

Protocols	Support for Micro mobility	Handover latency Micro mobility	Access Networks Support	Security	Legend
CMIP		Non	WLAN/WIMAX	+	+Stryig advantage
HMIP	++	138ms*	WLAN/WIMAX	-	-Acceptage -Draw -Strong drawback *not include network selection and luthentication **include authentication and manual network selection latency ***include re-authentication
HAWAII	++	150ms*	WLAN/WIMAX	1-	
Cellular IP	+	300ms*	WLAN/WIMAX/ Cellular Mobile Network	~	
PMIP	++	70ms***	WLAN/WIMAX/ Cellular Mobile Network	++	



#### 5.2 Issues of IPv6 Migrations

Due to low IP address space available for even preasing terminals there is a need of IPv6 in the near future to deliver the services. MTLMM is an IETF working group working in PMIPv6 [12, 13], Specification of PMIPv6 is still in the bfancy stage, there are several issues which has to be addressed to obtain the mobility solution. It was of Mobile IPv6 and PMIPv6 interactions, AAA support for PMIP, MPA discovery in the access networks, handover and route optimizations, Path Management and Failure Detection, Inter access handover support and multi homing scenario handover are still open in the WG. Using AAA mobility extensions and PMIPv6 supporting AAA extensions as proposed in the prchitecture, issues mentioned above are solved. As part of our future work we are developing adal stack PMIPv4 and PMIPv6 for mobility support in heterogeneous networks.

# 6 CONCLUSION

Post handover techniques are intended to reduce latency during roaming and handover in heterogeneries networks. As a part of this we have proposed security authentication and mobility management to potimize handover and roaming. Extending existing infrastructure such as AAA in this case is more officient than proposing new protocols and infrastructure. As a part of it security and mobility extensions are proposed. Using the security mechanisms we estimated the latency obtained in this method is far less than any conventional methods available in the literature. The authentication beying material created dynamically, by this way the theft presentational and security vulnerabilities are reduced. The mechanisms presented are applicable to WLAN, WIMAX and cellular networks and utilizing with RII architecture the solution provides the flexibility to over ate in any interworking scenarios of roaming and handover.

Proxy Mobile IP is a development of Mobile IP, where the registration is processed by the network entities. Hence, the Mobile Node does not require a Mobile IP stack to roam over the network without losing its IP address, so this can be applied to unchanged devices. Using this proposed mechanism, authentication and mobility management of users during the access is

performed in parallel; in this way, latency during the authentication and re-authentication is reduced. In this mechanism, using context management the control of users can be maintained according to the access networks. Fast and seamless handover is achieved in varius deployment and mobility scenarios using these mechanisms. Extending and upgrading existing networks can be performed efficiently, as no new hardware is added to the existing architectures. Multi homing scenarios, different interworking architectures of WLAN, WIMAX and 3G are addressed using the proposed mechanisms.

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